

## II. Forking and Lascar strong types in NIP theories

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### 1 Forking

#### Introduction

- Until recently I believed that Shelah's theory of forking and dividing was only meaningful in simple theories.
- In fact in theories with *NIP* it turns out to be very meaningful (as Shelah saw some time ago).
- In many examples of *NIP* theories where there is already a well-behaved notion of independence and dimension (such as *o*-minimal theories) forking can come from a lowering of "order of magnitude" within a given dimension (and is related to finding the right notion of "infinitesimal").
- Forking also gives rise to very nice analogues of stable group theory within suitable *NIP* theories (although this won't figure in my current talks).
- The level of this lecture and these notes will be notched up a bit compared to the introductory lecture.

#### Forking in general

- Definition 1.1.**
- (i) Let  $\Sigma(x)$  be a partial type (even over the big model), closed under finite conjunctions.  $\Sigma(x)$  is said to divide over  $A$  if there is  $\phi(x, b) \in \Sigma$  and an  $A$ -indiscernible sequence  $(b_i : i < \omega)$  of realizations of  $tp(b/A)$  such that  $\{\phi(x, b_i) : i < \omega\}$  is inconsistent.
  - (ii)  $\Sigma(x)$  is said to fork over  $A$  if  $\Sigma$  implies some finite disjunction of formulas (over  $\bar{M}$ ) each of which divides over  $A$ .

We begin with some trivialities (left to the reader) which also explain why the notion of forking, rather than just dividing, is introduced.

**Lemma 1.2.** (i)  $\Sigma(x)$  does not fork over  $A$  iff for any  $B$  containing  $\text{dom}(\Sigma)$ , (in particular for  $B = \bar{M}$ ),  $\Sigma$  extends to a complete type  $p(x)$  over  $B$  which does not divide over  $A$ .

(ii) If  $p(x)$  is a global complete type (or even a complete type over a sufficiently saturated model), and  $A$  is small, then  $p$  does not divide over  $A$  iff  $p$  does not fork over  $A$ .

- We first describe the situation when  $T$  is stable.
- Any partial type  $\Sigma(x)$  does not fork over its domain.
- Forking equals dividing in the sense that a partial type  $\Sigma(x)$  forks over  $A$  iff it divides over  $A$ .
- Let  $p(x) \in S(\bar{M})$  be a global type, and  $A$  a (small) subset of  $\bar{M}$ .
- Then  $p$  does not fork over  $A$  iff  $p$  is definable over  $\text{acl}^{eq}(A)$  iff  $p$  is almost finitely satisfiable in  $A$  (i.e. is finitely satisfiable in any model containing  $A$ ).
- Moreover  $p$  is the unique global type which extends  $p|_{\text{acl}^{eq}(A)}$  and does not fork over  $A$ .
- We also have the “algebraic” properties of (non)forking: symmetry, transitivity, existence of nonforking extensions,..., which are valid in any simple theory.

### Invariant types

**Definition 1.3.** Let  $p(x) \in S(\bar{M})$  (a global type) and  $A$  small.

- (i) We say that  $p$  does not split over  $A$  if  $p$  is  $\text{Aut}(\bar{M}/A)$ -invariant, equivalently whether or not  $\phi(x, b) \in p$  depends on  $tp(b/A)$ .
- (ii) We say that  $p$  does not Lascar-split over  $A$  if  $p$  is  $\text{Aut}f(\bar{M}/A)$ -invariant, equivalently whether or not  $\phi(x, b) \in p$  depends on  $Lstp(b/A)$ .
- (iii) We say (with some abuse of language) that  $p$  is invariant if  $p$  is  $\text{Aut}(\bar{M}/B)$ -invariant for some small  $B$ .
- (So if  $p$  does not Lascar split over a small set then  $p$  is invariant (why?).)

- If the global complete type  $p(x)$  is invariant then it has a kind of “infinite” defining schema (over a small set  $A$ ), namely for any  $\phi(x, y) \in L$ , there is a collection  $d_p\phi$  of complete  $y$ -types over  $A$  such that  $\phi(x, b) \in p$  iff  $tp(b/A) \in d_p\phi$ .
- So if  $B$  is any set containing  $A$  (and even containing  $\bar{M}$  in some saturated elementary extension of  $\bar{M}$ , we can apply this schema to  $B$  to obtain  $p|B$ .
- Suppose  $p(x), q(y)$  are global invariant types. Then by  $pq \in S_{xy}(\bar{M})$  we mean  $tp(a, b/\bar{M})$  where  $b$  realizes  $q$  and  $a$  realizes  $p|Mb$ . We can also form  $qp \in S_{xy}(\bar{M})$  as  $tp(a, b/\bar{M})$  where  $a$  realizes  $p$  and  $b$  realizes  $q|Ma$ .
- We let  $pq = qp$ .
- Let  $p(x)$  be an invariant global complete type, and  $x_1, x_2, \dots$  disjoint copies of the variable  $x$ .
- Then  $p^{(n)}(x_1, \dots, x_n)$  is defined inductively as  $p^{(n-1)}(x_1, \dots, x_{n-1})p(x_n)$ . And  $p^{(\omega)}(x_i)_i$  is the union of the  $p^{(n)}(x_1, \dots, x_n)$ .
- So  $p^{(\omega)}$  is  $tp(a_1, a_2, a_3, \dots/\bar{M})$  where  $a_1$  realizes  $p$ ,  $a_2$  realizes  $p|\bar{M}a_1$  etc.
- Any realization  $(a_i)_{i < \omega}$  of  $p^{(\omega)}(x_i)_i$  (with  $a_i$  corresponding to  $x_i$ , is easily seen to be an indiscernible sequence over  $\bar{M}$ .
- Assuming  $p$  to be  $Aut(\bar{M}/A)$ -invariant (or even  $Autf(\bar{M}/A)$ -invariant we may call any realization of  $p^{(\omega)}|A$  in  $\bar{M}$  a Morley sequence in  $p$  over  $A$ .

**Lemma 1.4.** • *i) Suppose that  $p(x) \in S(\bar{M})$  is  $Aut(\bar{M}/A)$ -invariant. Then so is  $p^{(\omega)}(x_i)_i$ .*

- *(ii) If  $p(x) \in S(\bar{M})$  is  $Autf(\bar{M}/A)$ -invariant, then so is  $p^{(\omega)}$ .*
- The proof is left as an exercise.
- Part (i) is easy.
- However part (ii) is a bit more tricky and uses Lemma 1.6 of Lecture I.

### Back to NIP

The following is always true (not requiring NIP).

**Lemma 1.5.** *If  $p(x) \in S(\bar{M})$  is  $Autf(\bar{M}/A)$ -invariant then  $p$  does not divide (fork) over  $A$ .*

*Proof.* If  $(b_i : i < \omega)$  is an  $A$ -indiscernible sequence then all the  $b_i$  have the same Lascar strong type over  $A$ , so if  $\phi(x, b_0) \in p$  then each  $\phi(x, b_i) \in p$  hence  $\{\phi(x, b_i) : i < \omega\}$  is consistent.  $\square$

The main point is that assuming NIP the converse to Lemma 1.5 holds.

**Theorem 1.6.** *Assume that  $T$  has NIP.*

- (i) *If  $p(x) \in S(\bar{M})$  does not fork over (small)  $A$ , then  $p$  is  $\text{Aut}(\bar{M}/A)$ -invariant (so with earlier terminology,  $p$  is invariant).*
- (ii) *In particular if  $M$  is a model, and  $p$  does not divide (fork) over  $M$  then  $p$  is  $\text{Aut}(\bar{M}/A)$ -invariant.*

*Proof.*

- Suppose  $p$  does not fork over  $A$ . By Lemma 1.6 of Lecture I it is enough to show that if  $\phi(x, y) \in L$  and  $(b_i : i < \omega)$  is  $A$ -indiscernible, then  $\phi(x, b_0) \in p$  iff  $\phi(x, b_1) \in p$ .
- If not then  $\phi(x, b_0) \Delta \phi(x, b_1) \in p(x)$ .
- But as  $(b_0 b_1, b_2 b_3, \dots)$  is also indiscernible over  $A$ , our assumptions imply that  $\{\phi(x, b_i) \Delta \phi(x, b_{i+1}) : i = 0, 2, \dots\}$  is consistent, contradicting Corollary 2.6 of Lecture I.

**Lemma 1.7.** (NIP) *Suppose  $p(x) \in S(\bar{M})$  is  $\text{Aut}(\bar{M}/A)$ -invariant. Let  $I$  be a Morley sequence in  $p$  over  $A$  (namely a realization of  $p^{(\omega)}|A$ ). Then  $I$  is  $A$ -special, and  $p = \text{Ev}(I/\bar{M})$ . In particular  $p$  is determined by the type of a Morley sequence in  $p$  over  $A$ .*

*Proof.*

- If  $I, J$  are realizations of  $p^{(\omega)}|A$  then clearly if  $I'$  is a realization of  $p^{(\omega)}|IJ$  then  $II'$  and  $IJ$  are  $A$ -indiscernible.
- Moreover for any  $B \supseteq A$  and realization  $I$  of  $p^{(\omega)}$ , if  $I'$  realizes  $p^{(\omega)}|BI$  then obviously  $\text{Ev}(I'/B) = p|B$ . This suffices.

Remark: We could refine the notion of an  $A$ -indiscernible sequence  $I$  being  $A$  special, by requiring that whenever  $I'$  and  $I$  have the same Lascar strong type over  $A$  then there is  $J$  etc. With a suitable definition of  $\text{Ev}(I/\bar{M})$  one could then write a version of Lemma 1.7 with “not forking over  $A$ ” in place of  $\text{Aut}(\bar{M}/A)$ -invariant.

**Corollary 1.8.** (NIP) *For any  $A$  and  $p(x) \in S(A)$ ,  $p$  has at most  $2^{(|A|+\omega)}$  global nonforking extensions.*

*Proof.* Exercise.

On the other hand one can prove that if  $T$  has the independence property then for any  $\lambda \geq |T|$  there is a model  $M$  of cardinality  $\lambda$  and a type  $p(x)$  over  $M$  with at least  $2^{2^\lambda}$  global nonforking extensions (in fact coheirs). So together with Corollary 1.8 this gives a nice characterization of theories with NIP.

- We can formulate the notion of a subset of (a sort in)  $\bar{M}$  being “Borel”.
- A set is “closed” if it is the solution set of a partial type over a small set.
- A set is Borel if it is in the  $\sigma$ -algebra generated by the closed sets.
- We could also call a set Borel over  $A$ , if it is in the  $\sigma$ -algebra generated by the solution sets of partial types over  $A$ .
- We call a set *strongly Borel* (over  $A$ ) if it is a finite Boolean combination of closed (over  $A$ ) sets.

**Theorem 1.9.** (*NIP*) *Let  $p(x) \in S(\bar{M})$  be  $\text{Aut}(\bar{M}/A)$ -invariant. Then  $p$  is strongly Borel definable over  $A$  in the sense that it has a strongly Borel over  $A$  defining schema, namely for any formula  $\phi(x, y)$  there is a strongly Borel set  $Y$  such that for any  $b$ ,  $\phi(x, b) \in p$  iff  $b \in Y$ .*

*Proof.* We use Lemma 2.5 of Lecture I.

- Let  $\phi(x, y) \in L$ , and  $N = n_\phi$  given by Lemma I.2.2(iii). Let  $Q(x_i)_i$  be  $p^{(\omega)}|A$  and  $Q_n(x_1, \dots, x_n)$  its restriction to  $x_1, \dots, x_n$ .
- Then  $\phi(x, b) \in p$  iff for some  $m \leq N$  there is  $(a_1, \dots, a_m)$  realizing  $Q_m$  with maximal possible successive alternations of truth values of  $\phi(x, b)$ , and moreover  $\models \phi(a_m, b)$ .
- I leave it as an exercise to express this as a strongly Borel over  $A$  condition.

## 2 Lascar strong types

### Invariant types and $KP$ -strong types

We begin with a result valid at the level of “invariant types”, which then applies to the *NIP* context.

**Lemma 2.1.** *Let  $p(x) \in S(\bar{M}/A)$  be  $\text{Aut}(\bar{M}/A)$ -invariant. Let  $c, d$  be realizations of  $p$ . Then  $\text{Lstp}(c/A) = \text{Lstp}(d/A)$  iff there is some infinite sequence  $\bar{a}$  such that both  $(c\bar{a})$  and  $(d\bar{a})$  realize  $p^{(\omega)}|A$ .*

*Proof.* By Lemma I.1.6, RHS implies LHS.

Now assume RHS. After possibly moving  $p$  by an  $A$ -automorphism, we may assume that  $p(x)$  implies  $\text{Lstp}(c/A)$ . Let  $(a_0, a_1, \dots)$  realize  $p^{(\omega)}$ . Then  $(a_1, a_2, \dots)$  also realizes  $p^{(\omega)}$ . By Lemma 1.4(ii),  $p^{(\omega)}$  is  $\text{Aut}(\bar{M}/A)$ -invariant. As  $\text{Lstp}(c/A) = \text{Lstp}(a_0/A) = \text{Lstp}(d/A)$  we conclude that  $\text{tp}(c, a_1, a_2, \dots/A) = \text{tp}(a_0, a_1, a_2, \dots/A) = \text{tp}(d, a_1, a_2, \dots/A)$ , completing the proof.

**Corollary 2.2.** *Suppose  $p(x) \in S(\bar{M})$  is  $\text{Aut}(\bar{M}/A)$ -invariant. Then for any realization  $a$  of  $p|A$ ,  $KPstp(a/A)$  equals  $Lstp(a/A)$*

*Proof.* Lemma 2.1 implies that on realizations of  $p|A$ , the relation  $Lstp(x_1/A) = Lstp(x_2/A)$  is type-definable over  $A$ , which suffices.

Restatement for *NIP* theories.

**Corollary 2.3.** *(NIP) Suppose  $p(x) \in S(A)$  does not fork over  $A$ . Then on realizations of  $p$ ,  $KPstp$  over  $A$  equals  $Lstp$  over  $A$ .*

*Proof.* Exercise using earlier results and propositional calculus.

**Corollary 2.4.** *(NIP) Suppose that  $T$  is 1-sorted, and that for every  $A$ , every complete 1-type over  $A$  does not fork over  $A$  (equivalently has a global nondividing extension). Then over any set  $A$ ,  $KPstp = Lstp$ .*

*Proof.* Use a result of Shelah: assuming *NIP* if  $tp(a/B)$  does not fork over  $A$  and  $tp(b/Ba)$  does not fork over  $A$ , then  $tp(a, b/B)$  does not fork over  $A$ , to conclude that any complete type over any  $A$  does not fork over  $A$ . Then use the previous Corollary.

### **Forking and $bdd(A)$ -invariance**

In this final minisection we aim to prove the following strengthening of Theorem 1.6.

**Theorem 2.5.** *(NIP) Let  $p(x) \in S(\bar{M})$  and  $A$  a small set. Then  $p$  does not fork over  $A$  iff  $p$  is  $bdd^{heq}(A)$ -invariant (namely  $p$  is fixed by all automorphisms which fix all  $KP$ -strong types over  $A$ ).*

Sketch of proof.

- This really falls out of our previous results.
- (i) Note that (by Lemma 1.4 for example)  $p^{(\omega)} \in S_\omega(\bar{M})$  does not fork over  $A$ .
- (ii) So by (i) and Corollary 2.3 if  $I$  realizes  $p^{(\omega)}|A$  then  $KPstp(I/A)$  implies  $Lstp(I/A)$ .
- Let  $Q(x_i)_i = p^{(\omega)}|bdd^{heq}(A)$  (a  $KP$ -strong type over  $A$ ).
- (iii) So (ii) says that  $Q$  implies a Lascar strong type over  $A$ .

- *Claim (iv).*  $Q$  is  $bdd^{heq}$ -special, more precisely  $tp(I/bdd^{heq}(A))$  is  $bdd^{heq}(A)$ -special for some (any)  $I$  realizing  $Q$ .
- *Proof of Claim (iv).* Let  $I_1, I_2$  realize  $Q$ , and let  $J$  realize  $p^{(\omega)}$  (in a saturated model containing  $\bar{M}$ ).
- Then  $Lstp(I_1/A) = Lstp(I_2/A) = Lstp(J/A)$ , hence as  $p^{(\omega)}$  is  $Aut(\bar{M}/A)$ -invariant, both  $I_1J$  and  $I_2J$  are  $A$ -indiscernible. End of proof of Claim (iv).
- Deduce as in Lemma 1.7 that  $p = Ev(I/\bar{M})$  for any realization  $I$  of  $Q$ .
- Hence  $p$  is  $Aut(\bar{M}/bdd^{heq}(A))$ -invariant as required.

**Corollary 2.6.** (*NIP*) Suppose  $p(x) \in S(\bar{M})$  does not fork over  $A$ . Then  $p$  is strongly Borel definable over  $bdd^{heq}(A)$ .